



Resilient and Fast Persistent Container Storage Leveraging Linux's Storage Functionalities

Philipp Reisner, CEO LINBIT

Presenter: Robert Altnoeder

LINBIT - the company behind it



COMPANY OVERVIEW

- Developer of DRBD
- 100% founder owned
- Offices in Europe and US
- Team of 30 highly experienced Linux experts
- Partner in Japan



TECHNOLOGY OVERVIEW



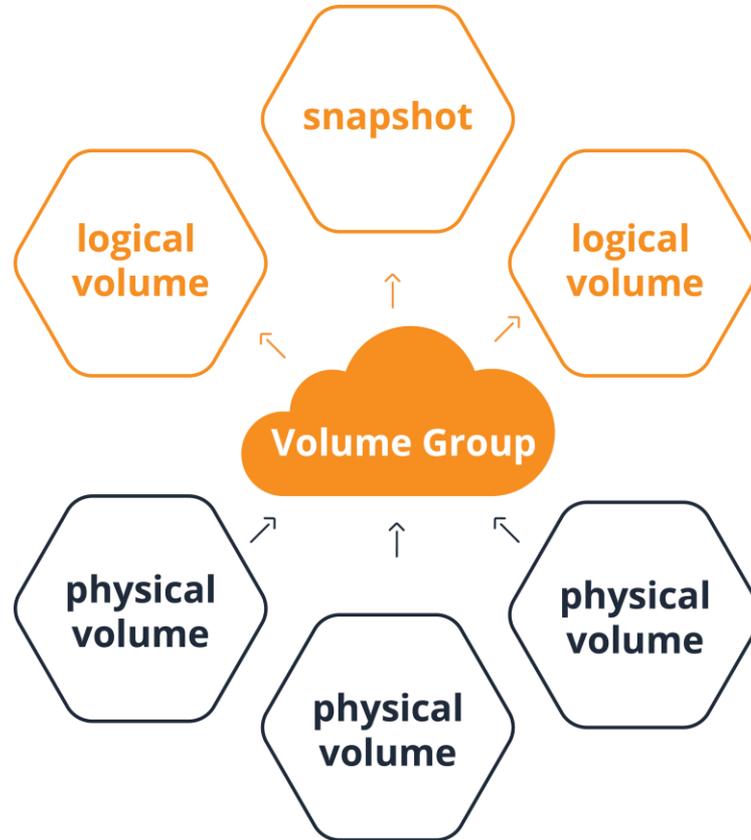
REFERENCES



Linux Storage Gems

LVM, RAID, SSD cache tiers, deduplication, targets & initiators

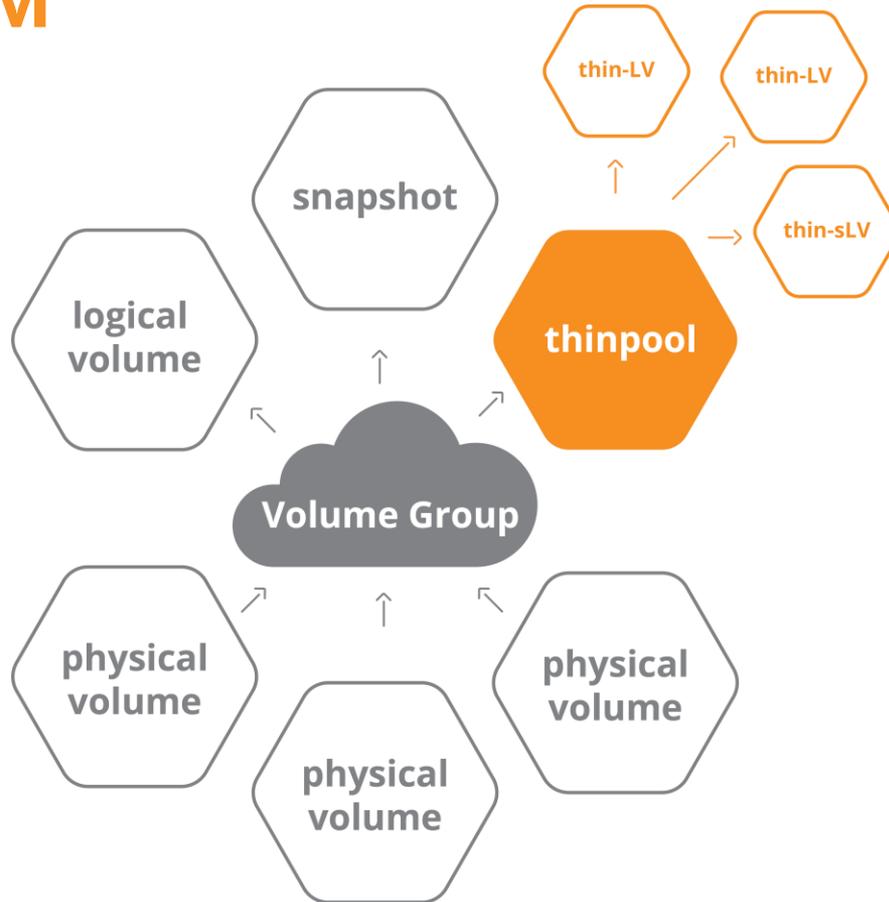
Linux's LVM



Linux's LVM

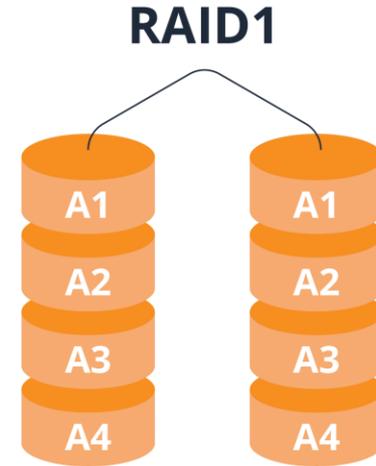
- based on device mapper
- original objects
 - PVs, VGs, LVs, snapshots
 - LVs can scatter over PVs in multiple segments
- thinlv
 - thinpools = LVs
 - thin LVs live in thinpools
 - multiple snapshots became efficient!

Linux's LVM



Linux's RAID

- original MD code
 - `mdadm` command
 - Raid Levels: 0,1,4,5,6,10
- Now available in LVM as well
 - device mapper interface for MD code
 - do not call it 'dmraid'; that is software for hardware fake-raid
 - `lvcreate --type raid6 --size 100G VG_name`



SSD cache for HDD

- dm-cache
 - device mapper module
 - accessible via LVM tools
- bcache
 - generic Linux block device
 - slightly ahead in the performance game



Linux's DeDupe

- Virtual Data Optimizer (VDO) since RHEL 7.5
 - Red hat acquired Permabit and is GPLing VDO
- Linux upstreaming is in preparation
- in-line data deduplication
- kernel part is a device mapper module
- indexing service runs in user-space
- async or synchronous writeback
- Recommended to be used below LVM

Linux's targets & initiators

- Open-ISCSI initiator
- letd, STGT, SCST
 - mostly historical
- **LIO**
 - iSCSI, iSER, SRP, FC, FCoE
 - SCSI pass through, block IO, file IO, user-specific-IO
- NVMe-OF
 - target & initiator



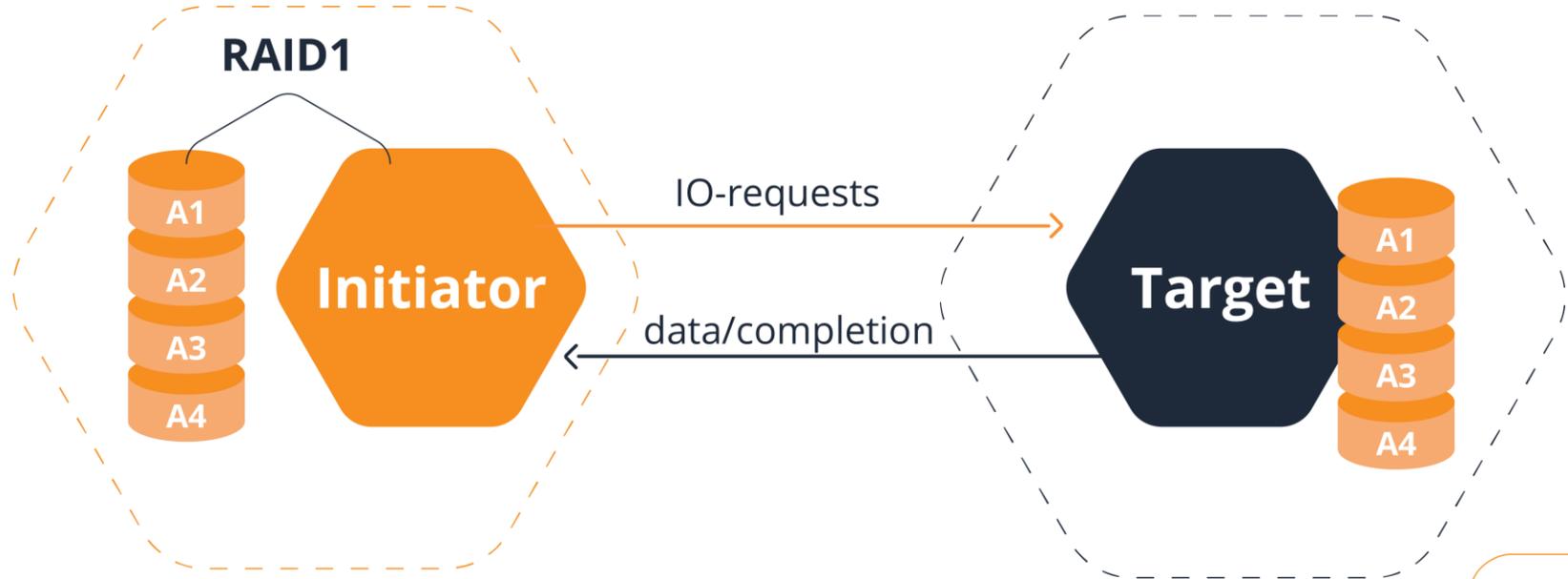
ZFS on Linux

- Ubuntu eco-system only
- has its own
 - logic volume manager (zVols)
 - thin provisioning
 - RAID (RAIDz)
 - caching for SSDs (ZIL, SLOG)
 - and a file system!

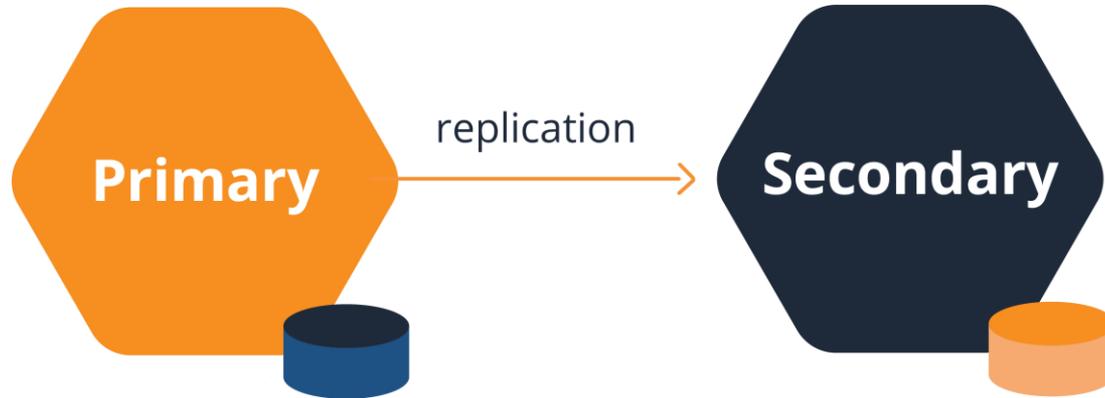
DRoBD

Put in simplest form

DRBD – think of it as ...

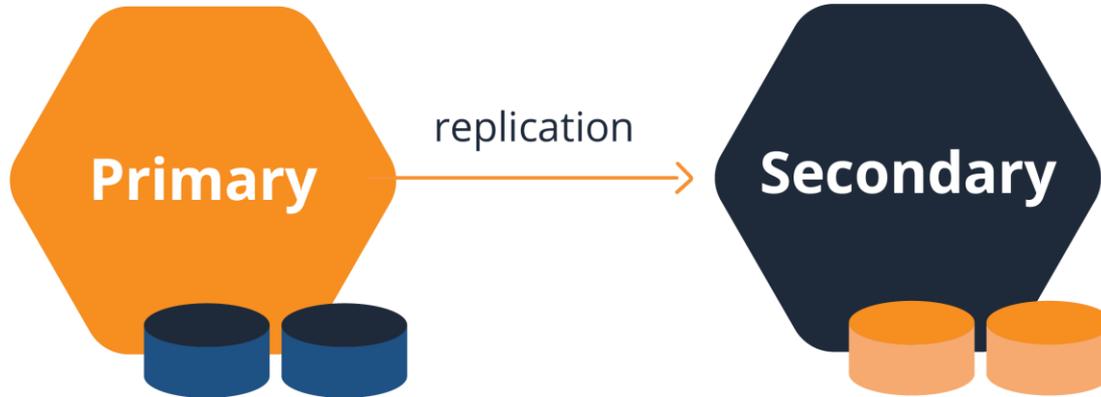


DRBD Roles: Primary & Secondary



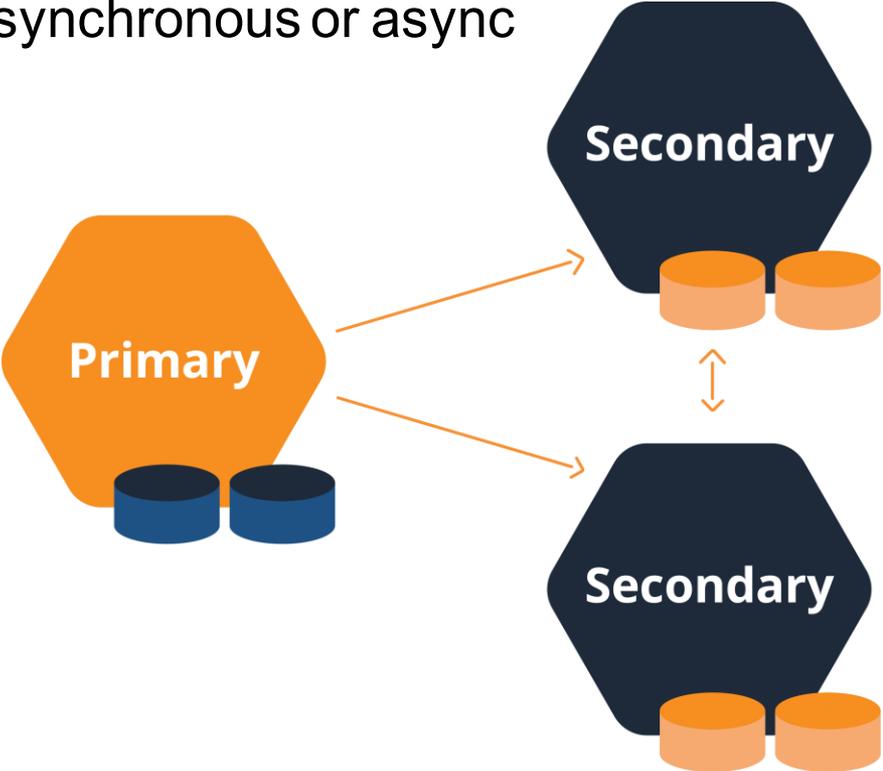
DRBD – multiple Volumes

- consistency group



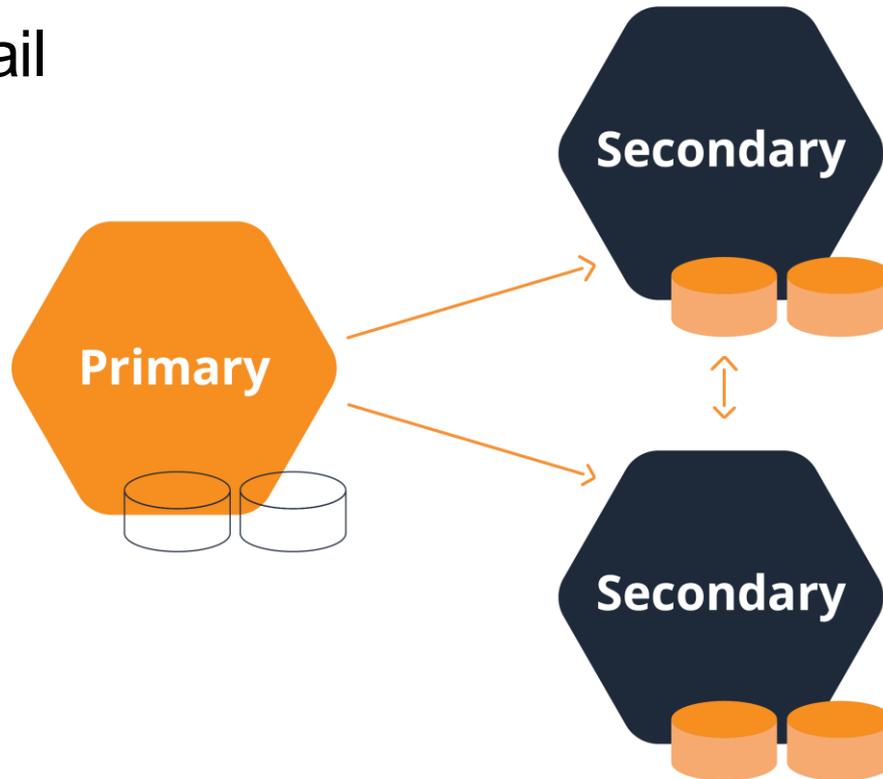
DRBD – up to 32 replicas

- each may be synchronous or async



DRBD – Diskless nodes

- intentional diskless (no change tracking bitmap)
- disks can fail



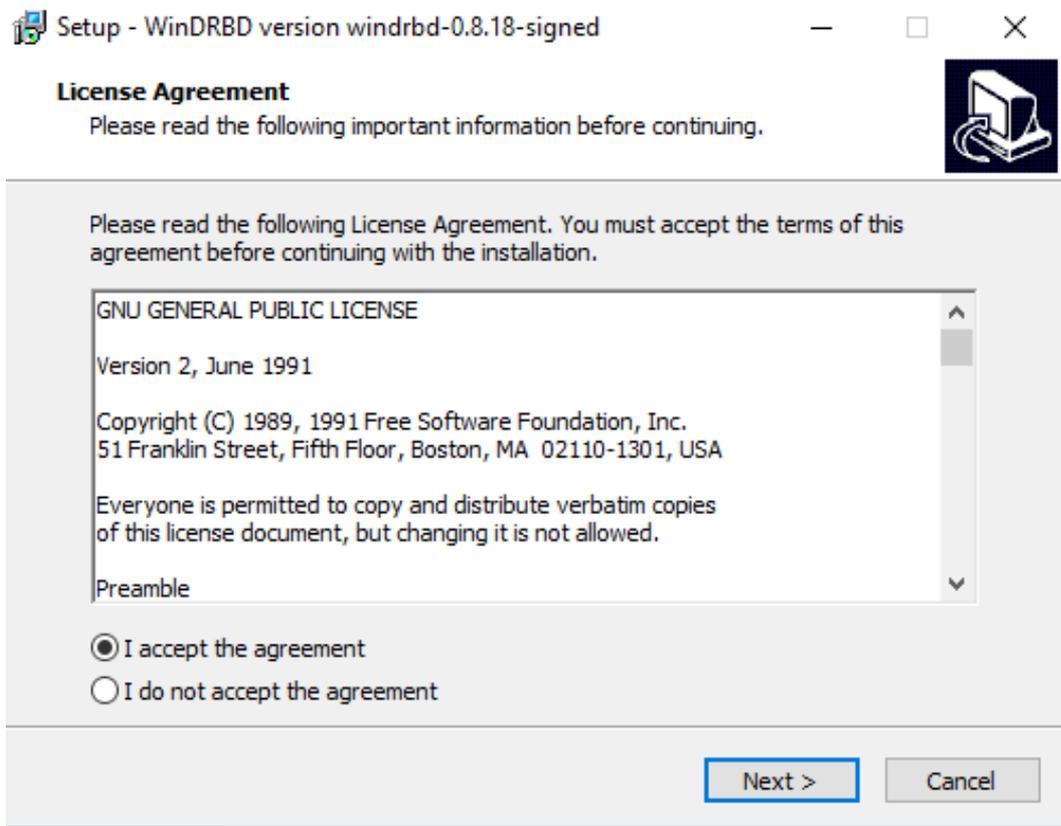
DRBD - more about

- a node knows the version of the data it exposes
- automatic partial resync after connection outage
- checksum-based verify & resync
- split brain detection & resolution policies
- fencing
- quorum
- multiple resources per node possible (1000s)
- dual Primary for live migration of VMs only!

DRBD Roadmap

- performance optimizations
 - meta-data on PMEM/NVDIMMS
 - zero copy receive on diskless (RDMA-transport)
 - no context switch send (RDMA & TCP transport)
 - Improve resync speed
- Eurostars grant: DRBD4Cloud
 - erasure coding (2019)
- Long distance replication
 - send data once over long distance to multiple replicas

WinDRBD



- in public beta
 - <https://www.linbit.com/en/drbd-community/drbd-download/>
- Windows 7sp1, Windows 10, Windows Server 2016
- wire protocol compatible to Linux version
- driver tracks Linux version with one day release offset
- WinDRBD user level tools are merged into upstream

WinDRBD ROADMAP 2019



- fix multiple connections (Februar)
- add auto-promote (March)
- enable WinDRBD for boot and drive C: (March, April)
- review/rework spinlock & RCU primitives (May)
- POCs with customers (starting in July)



The combination is more than the sum of its parts

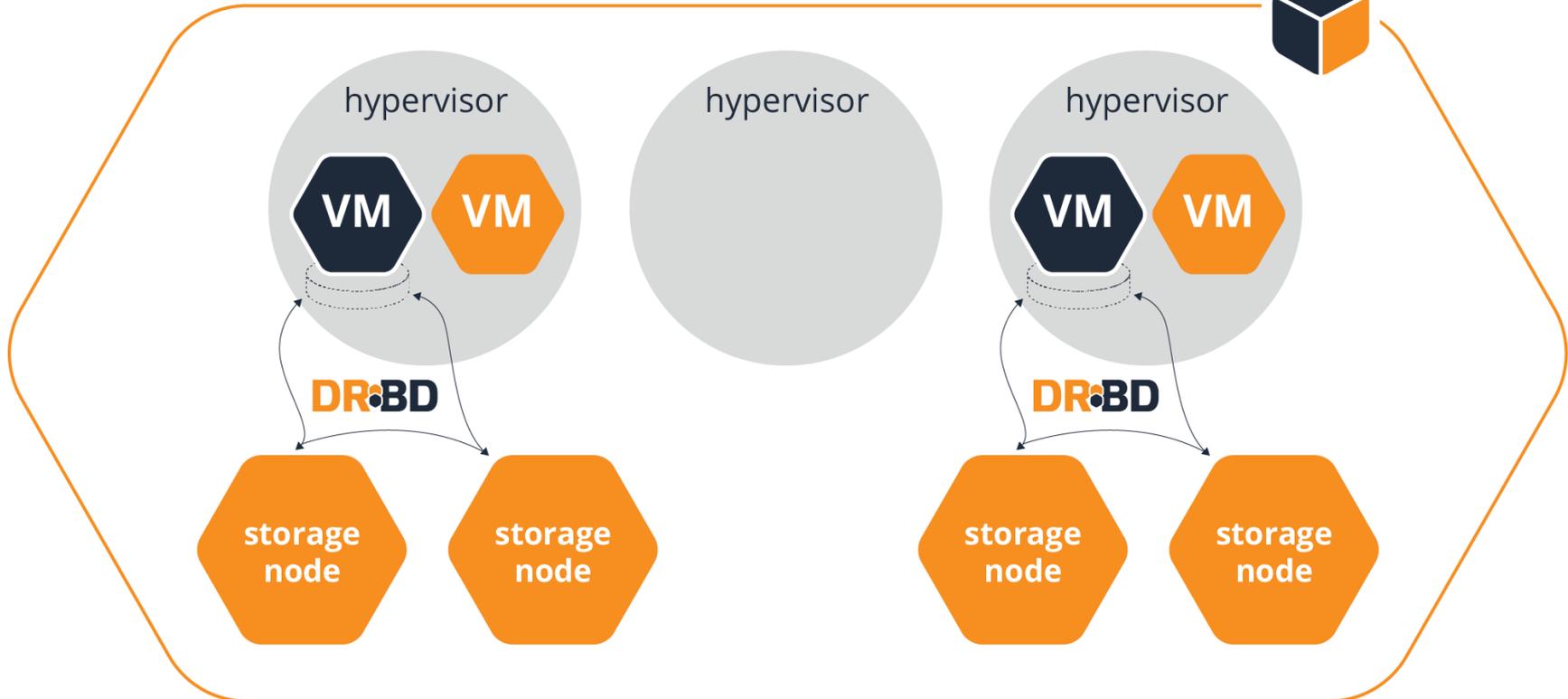
LINSTOR - goals

- storage built from generic (x86) nodes
- for SDS consumers (K8s, OpenStack, OpenNebula)
- building on existing Linux storage components
- multiple tenants possible
- deployment architectures
 - distinct storage nodes
 - hyperconverged with hypervisors / container hosts
- LVM, thin LVM or ZFS for volume management (stratis later)
- **Open Source, GPL**

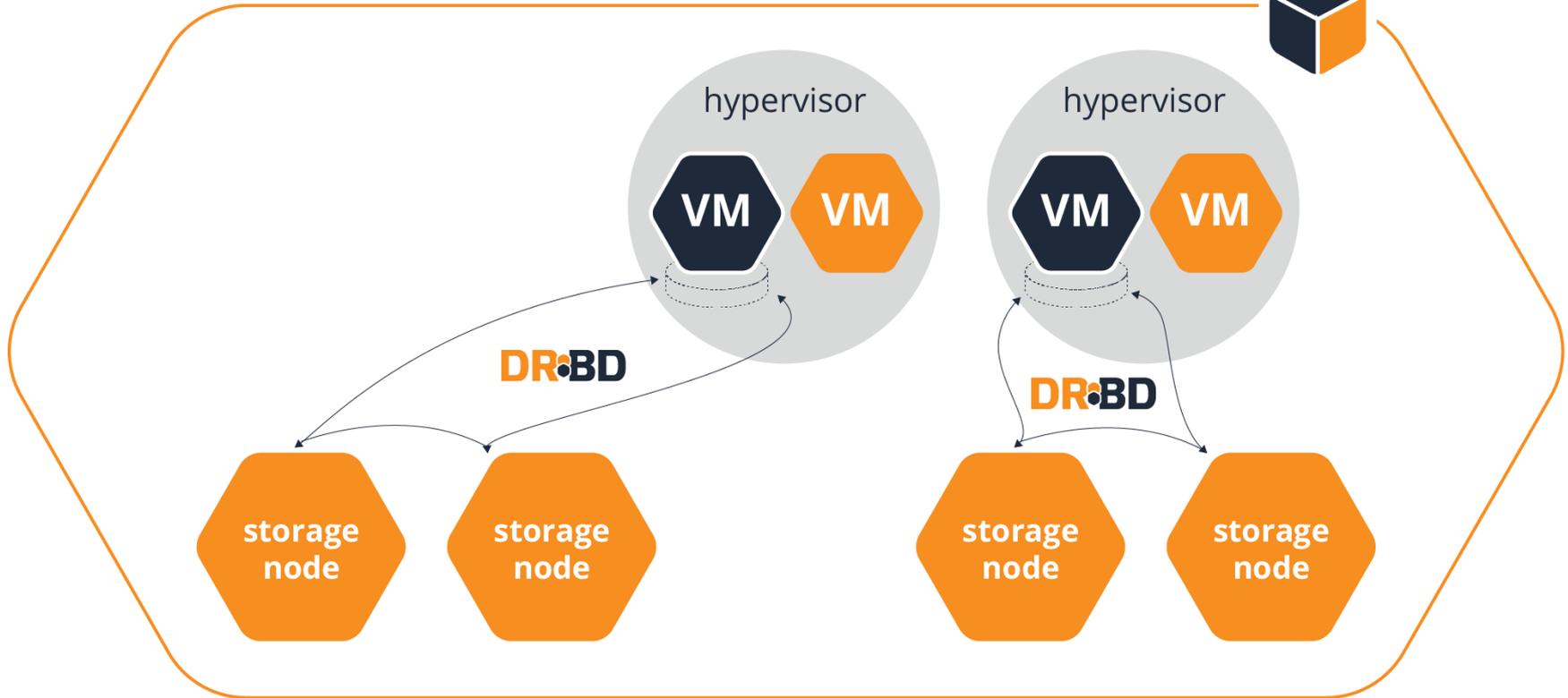


Examples

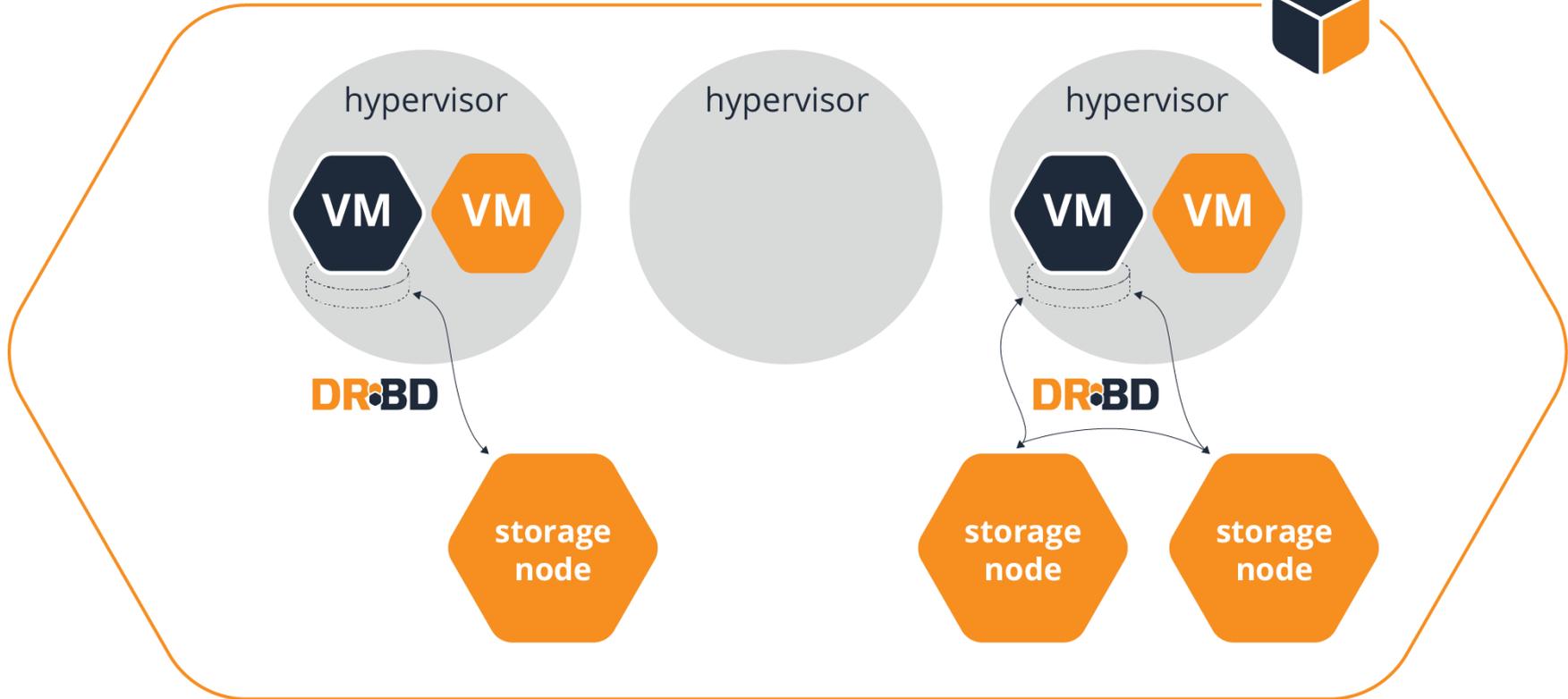
LINSTOR – disaggregated stack



LINSTOR / failed Hypervisor



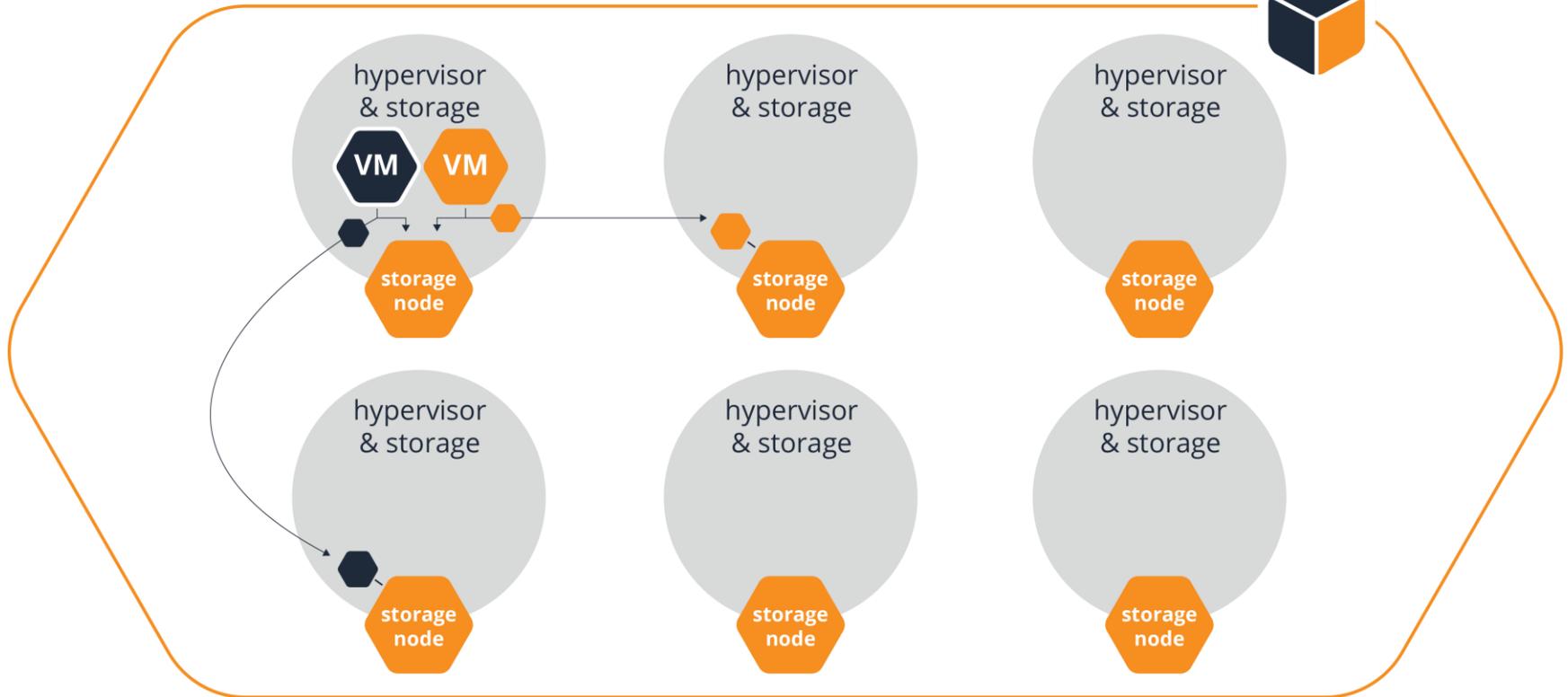
LINSTOR / failed storage node



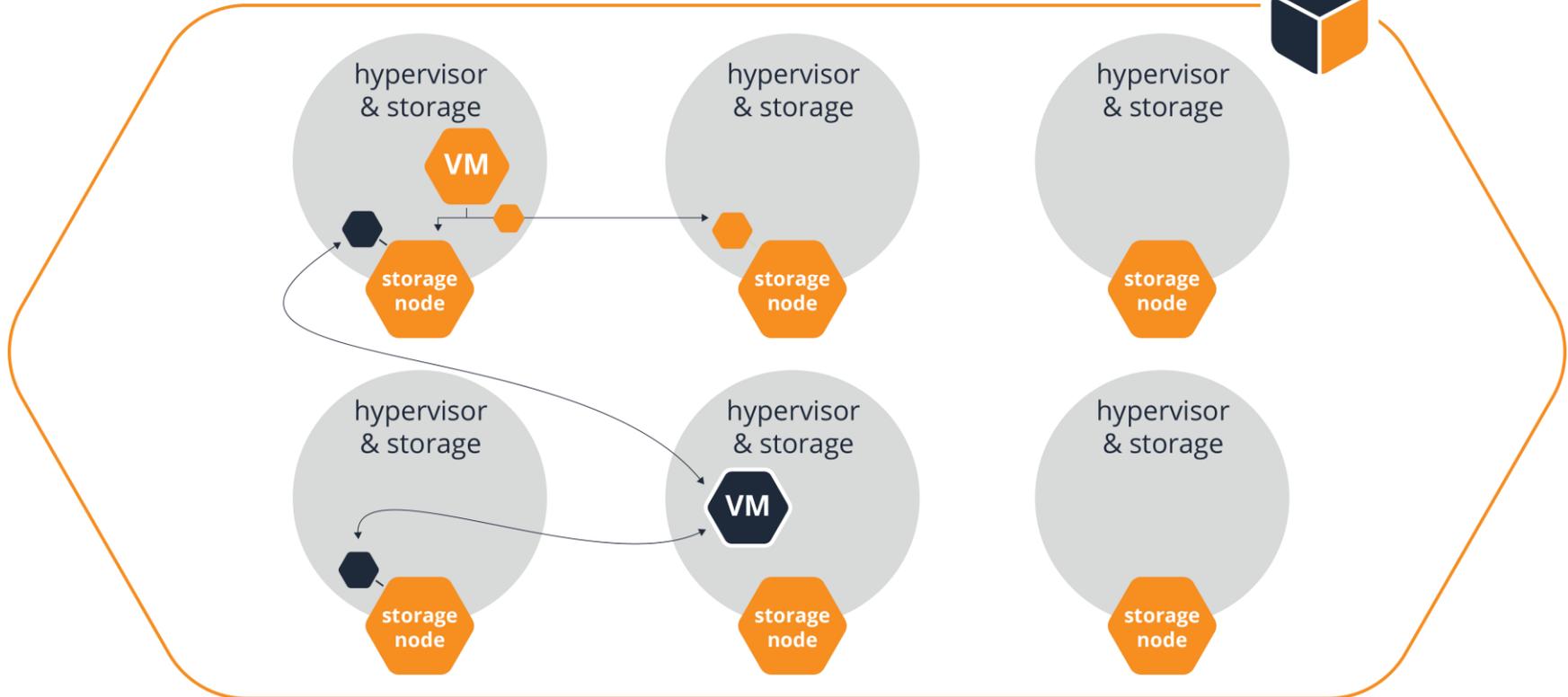
LINSTOR - Hyperconverged

LINBIT

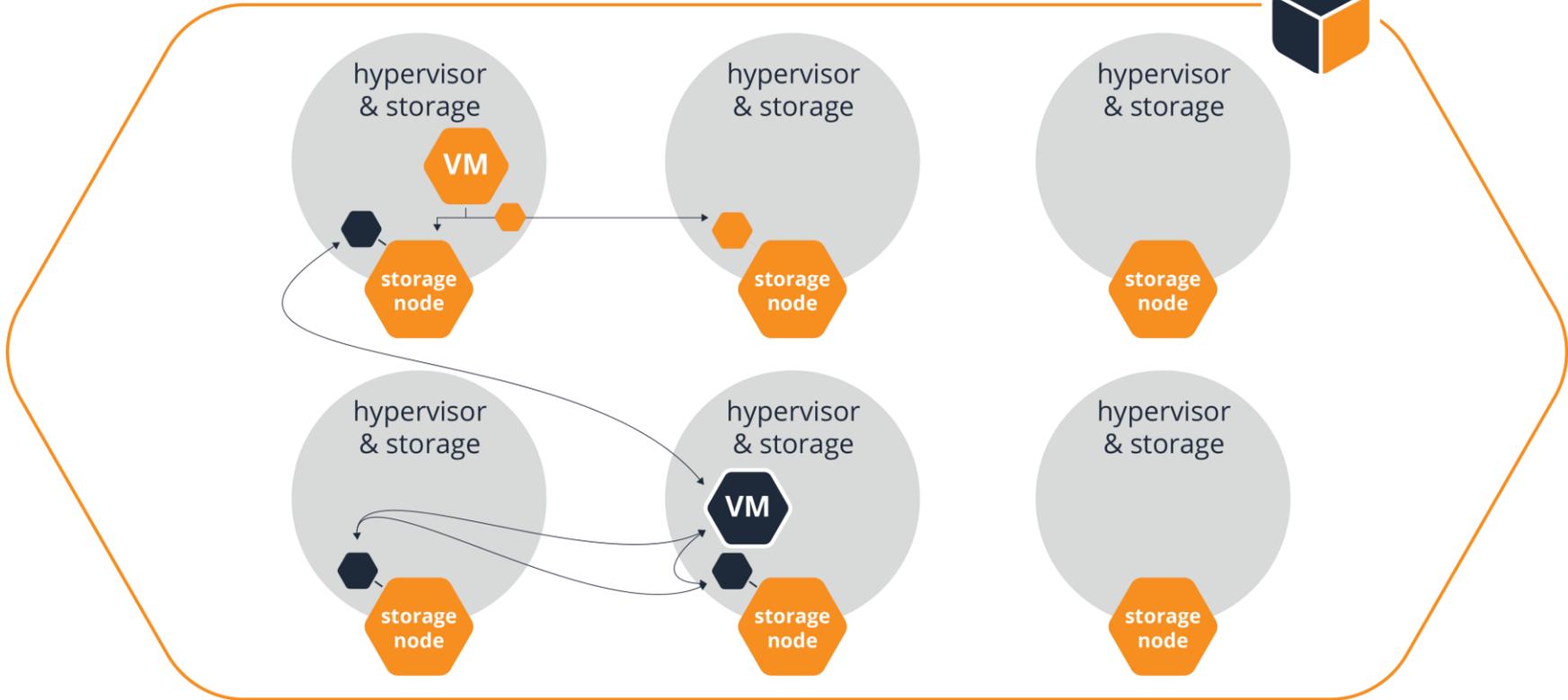
LINSTOR



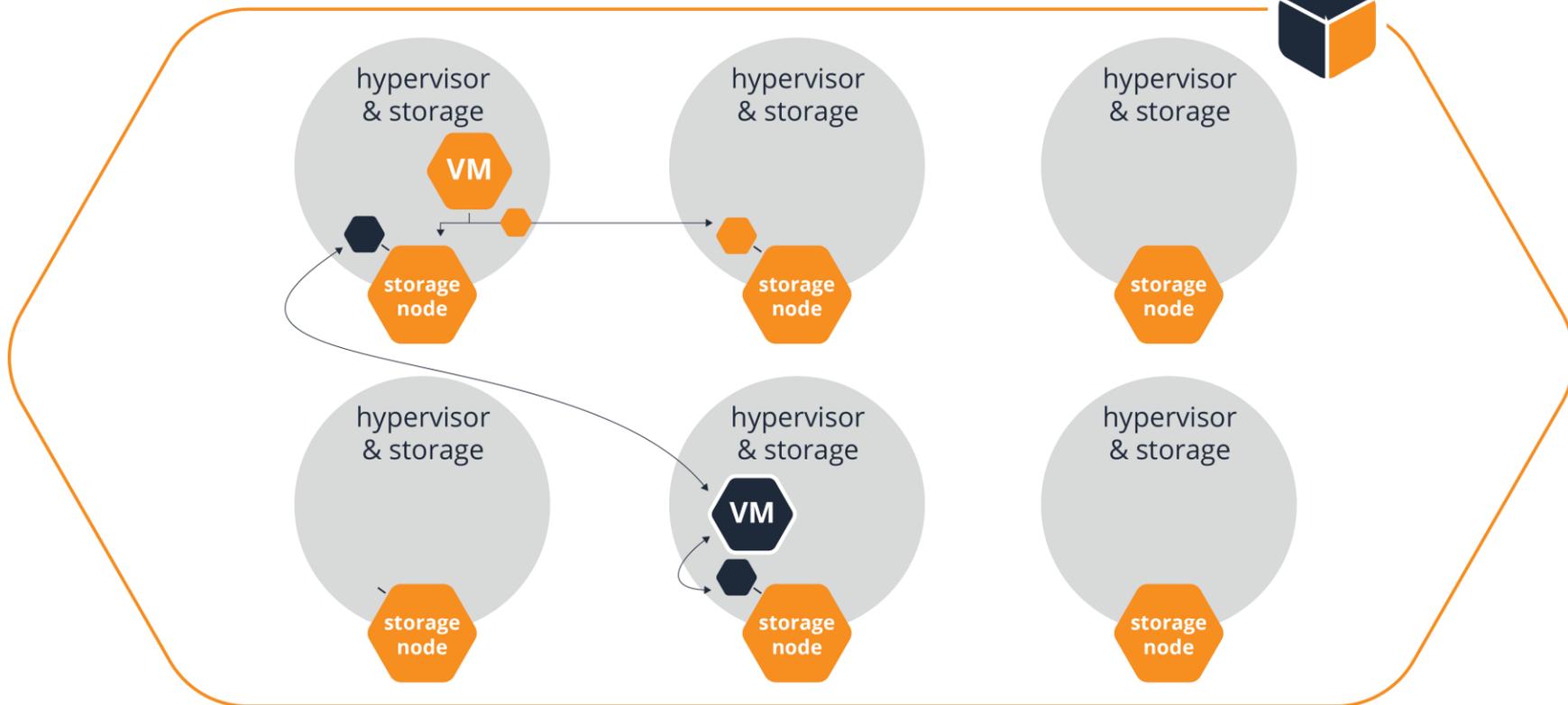
LINSTOR - VM migrated



LINSTOR - add local replica

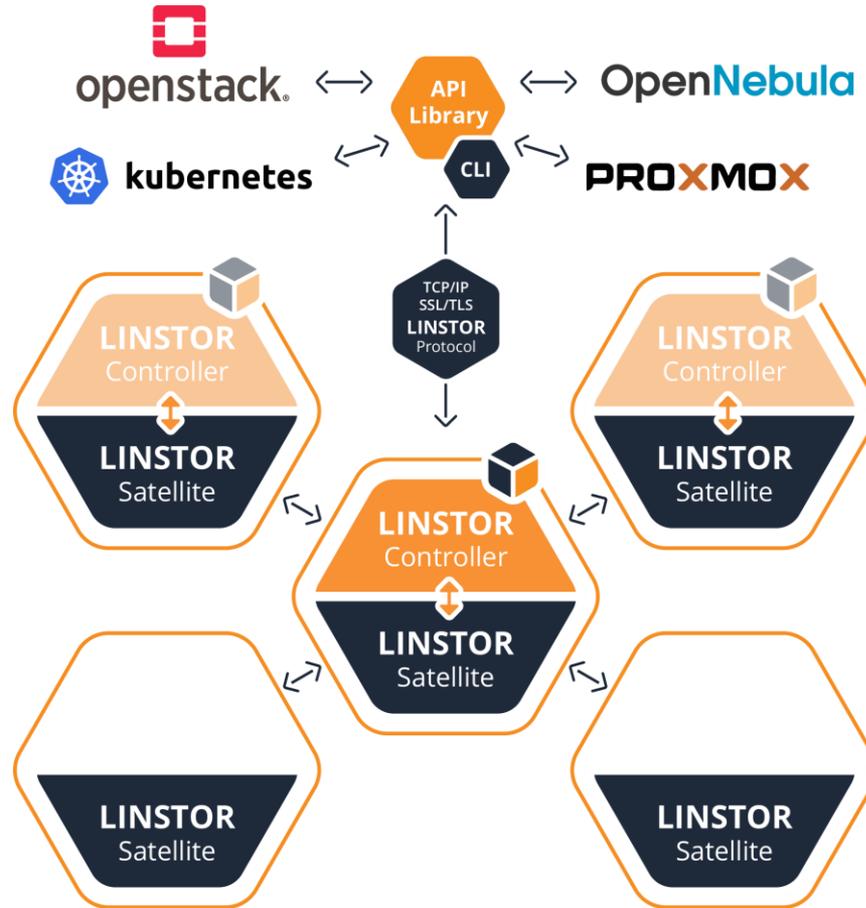


LINSTOR - remove 3rd copy



LIN  **STOR**

Architecture and functions



LINSTOR data placement

- arbitrary tags on nodes
 - require placement on equal/different/named tag values
- prohibit placements with named existing volumes
 - different failure domains for related volumes

Example policy

3 way redundant, where two copies are in the same rack but in different fire compartments (synchronous) and a 3rd replica in a different site (asynchronous)

Example tags

rack = number
room = number
site = city

LINSTOR network path selection

- a storage pool may prefer a NIC
 - express NUMA relation of NVMe devices and NICs
- DRBD's multi-pathing supported
 - load balancing with the RDMA transport
 - fail-over only with the TCP transport

LINSTOR connectors



- Kubernetes
 - FlexVolume & External Provisioner
 - CSI (Docker Swarm, Mesos)



- OpenStack/Cinder
 - since Stein release (April 2019)



- OpenNebula



- Proxmox VE



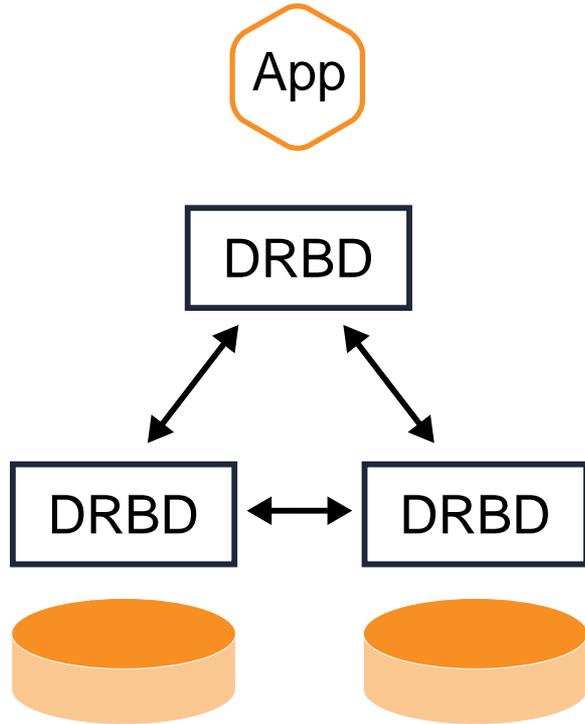
- XenServer / XCP-ng



LINSTOR Roadmap

- generalize LINSTOR for other IO stacks
 - MD-Raid & NVMe-oF
 - optional HW discovery & VG automatic VG creation
 - bcache & deduplication
- Linux NVMe-oF initiator & targets
- GUI based on REST-API
- auto-placement policies as LINSTOR objects
- LINSTOR & WinDRBD (?)

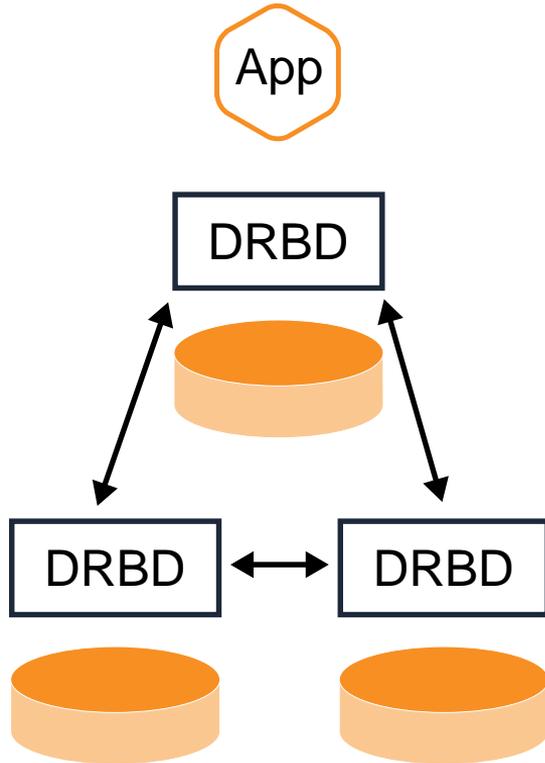
LINSTOR Storage Stacks



- Disaggregated Storage
- Classic enterprise workloads
 - Data bases
 - Message queues
- Typical Orchestrators
 - OpenStack, OpenNebula
 - Kubernetes
- Flexible redundancy (1-n)
- HDDs, SSDs, NVMe SSDs



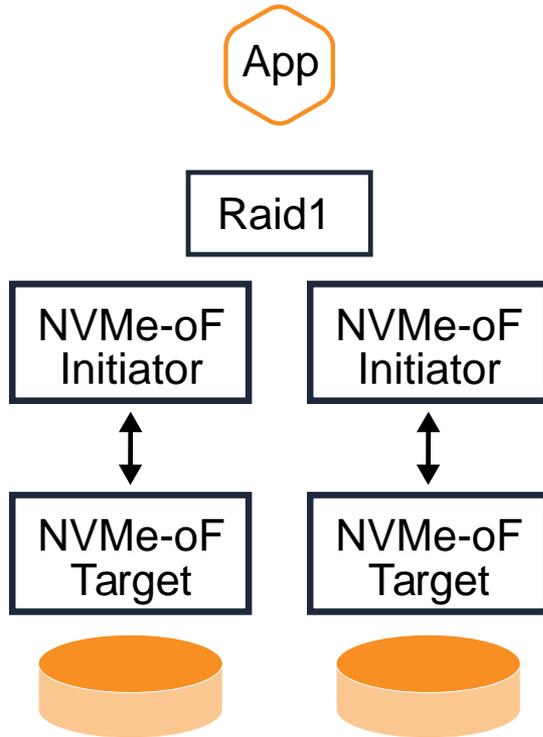
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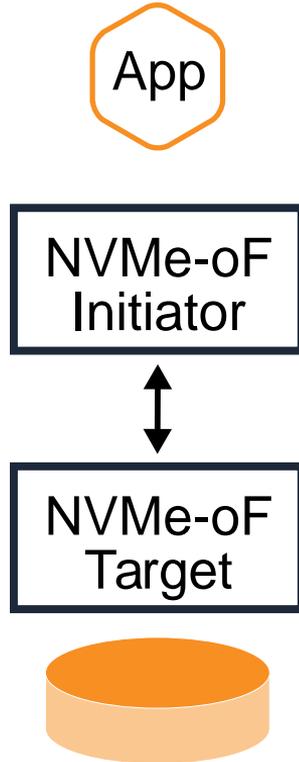
LINSTOR Storage Stacks



- Disaggregated
- Classic enterprise workloads
 - Data bases
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- Typical Orchestrators
 - OpenStack, OpenNebula
 - Kubernetes
- NVMe SSDs, SSDs



LINSTOR Storage Stacks



- Disaggregated
- Cloud native workload
 - Ephemeral storage
- Typical Orchestrator
 - Kubernetes
- Application handles redundancy
- Best suited for NVMe SSDs



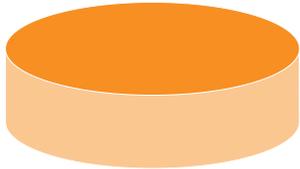
LINSTOR Storage Stacks



- Hyperconverged
- Cloud native workload
 - Ephemeral storage
 - PMEM optimized data base
- Typical Orchestrator
 - Kubernetes
- Application handles redundancy
- PMEM, NVDIMMs



LINSTOR Slicing Storage



- LVM or ZFS
- Thick – pre allocated
 - Best performance
 - Less features
- Thin – allocated on demand
 - Overprovisioning possible
 - Many snapshots possible
- Optional
 - Encryption on top
 - Deduplication below



Case study - intel



Intel® Rack Scale Design (Intel® **RSD**) is an industry-wide architecture for disaggregated, composable infrastructure that fundamentally changes the way a data center is built, managed, and expanded over time.

LINBIT working together with Intel

LINSTOR is a storage orchestration technology that brings storage from generic Linux servers and SNIA Swordfish enabled targets to containerized workloads as persistent storage. LINBIT is working with Intel to develop a Data Management Platform that includes a storage backend based on LINBIT's software. LINBIT adds support for the SNIA Swordfish API and NVMe-oF to LINSTOR.



Thank you

<https://www.linbit.com>